

Epistemic irrelevance in credal networks

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Before we start



Precise probability models

Mass functions and expectations

Assume we are **uncertain** about:

- the value or a variable X
- in a set of possible values \mathcal{X} .

This is usually modelled by a **probability mass function** p on \mathcal{X} :

$$p(x) \geq 0 \text{ and } \sum_{x \in \mathcal{X}} p(x) = 1;$$

With p we can associate an **expectation operator** E_p :

$$E_p(f) := \sum_{x \in \mathcal{X}} p(x)f(x) \text{ where } f: \mathcal{X} \rightarrow \mathbb{R}.$$

If $A \subseteq \mathcal{X}$ is an **event**, then its **probability** is given by

$$P_p(A) = \sum_{x \in A} p(x) = E_p(I_A).$$

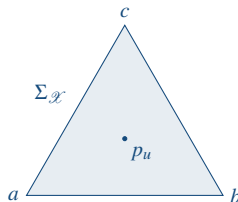
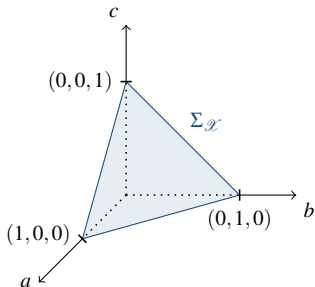


Precise probability models

The simplex of all probability mass functions

Consider the **simplex** $\Sigma_{\mathcal{X}}$ of all mass functions on \mathcal{X} :

$$\Sigma_{\mathcal{X}} := \left\{ p \in \mathbb{R}_+^{\mathcal{X}} : \sum_{x \in \mathcal{X}} p(x) = 1 \right\}.$$

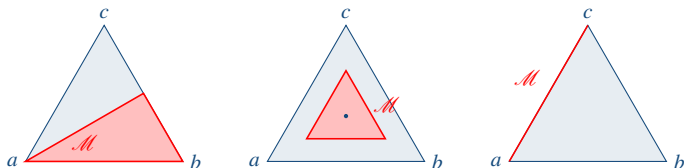


Imprecise probability models

Credal sets

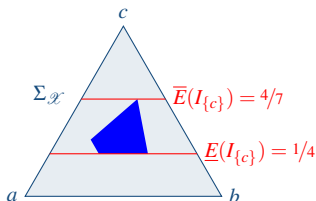
Definition

A credal set \mathcal{M} is a **convex closed** subset of $\Sigma_{\mathcal{X}}$.



Imprecise probability models

Lower and upper expectations



Equivalent model

Consider the set $\mathcal{L}(\mathcal{X}) = \mathbb{R}^{\mathcal{X}}$ of all real-valued maps on \mathcal{X} . We define two real functionals on $\mathcal{L}(\mathcal{X})$: for all $f: \mathcal{X} \rightarrow \mathbb{R}$

$$\underline{E}_{\mathcal{M}}(f) = \min \{E_p(f) : p \in \mathcal{M}\} \text{ lower expectation}$$

$$\bar{E}_{\mathcal{M}}(f) = \max \{E_p(f) : p \in \mathcal{M}\} \text{ upper expectation.}$$

Observe that [conjugacy]: $\bar{E}_{\mathcal{M}}(f) = -\underline{E}_{\mathcal{M}}(-f)$.

Imprecise probability models

Basic properties of upper expectations

Definition

We call a real functional \bar{E} on $\mathcal{L}(\mathcal{X})$ an **upper expectation** if it satisfies the following properties:

For all f and g in $\mathcal{L}(\mathcal{X})$ and all real $\lambda \geq 0$:

- 1 $\bar{E}(f) \leq \max f$ [**boundedness**];
- 2 $\bar{E}(f + g) \leq \bar{E}(f) + \bar{E}(g)$ [**sub-additivity**];
- 3 $\bar{E}(\lambda f) = \lambda \bar{E}(f)$ [**non-negative homogeneity**].

Theorem (Other properties)

Let \bar{E} be an upper expectation, with conjugate **lower expectation** \underline{E} . Then for all real numbers μ and all f and g in $\mathcal{L}(\mathcal{X})$:

- 1 $\underline{E}(f) \leq \bar{E}(f)$;
- 2 $\underline{E}(f) + \underline{E}(g) \leq \underline{E}(f + g) \leq \underline{E}(f) + \bar{E}(g) \leq \bar{E}(f + g) \leq \bar{E}(f) + \bar{E}(g)$;
- 3 $\bar{E}(f + \mu) = \bar{E}(f) + \mu$;
- 4 $\bar{E}(|f|) \geq |\underline{E}(f)|$ **and** $\bar{E}(|f|) \geq |\bar{E}(f)|$.

Imprecise probability models

Lower Envelope Theorem

Theorem (Lower Envelope Theorem)

A real functional \bar{E} is an upper expectation if and only if it is the upper envelope of some credal set \mathcal{M} .

Proof.

Use $\mathcal{M} = \{p \in \Sigma_{\mathcal{X}} : (\forall f \in \mathcal{L}(\mathcal{X}))(E_p(f) \leq \bar{E}(f))\}$. □

Types of independence

Three possible definitions

Epistemic irrelevance

X_2 is epistemically irrelevant to X_1 , conditional on X_3 :

$$\bar{E}(f(X_1)|X_2, X_3) = \bar{E}(f(X_1)|X_3)$$

Epistemic independence

X_1 and X_2 are epistemically independent, conditional on X_3 :

$$\bar{E}(f(X_1)|X_2, X_3) = \bar{E}(f(X_1)|X_3) \text{ and } \bar{E}(g(X_2)|X_1, X_3) = \bar{E}(g(X_2)|X_3)$$

Strong independence

Model $\bar{E}(h(X_1, X_2)|X_3)$ is an upper envelope of precise independent models

Discrete-time uncertain processes

Precise probability trees

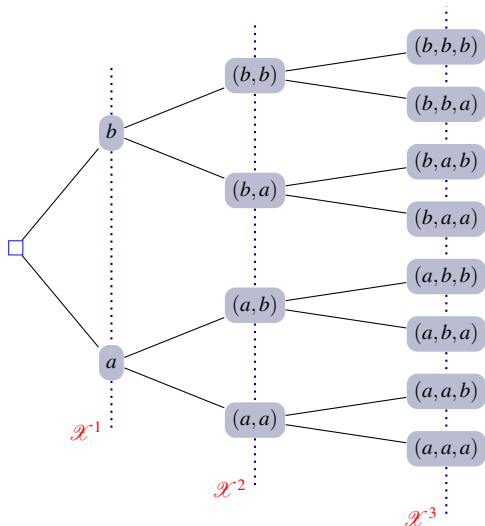
We consider an **uncertain process** with variables $X_1, X_2, \dots, X_n, \dots$ assuming values in a finite set of **states** \mathcal{X} .

This leads to a **standard event tree** with nodes

$$s = (x_1, x_2, \dots, x_n), \quad x_k \in \mathcal{X}, \quad n \geq 0.$$

Discrete-time uncertain processes

Precise probability trees



Discrete-time uncertain processes

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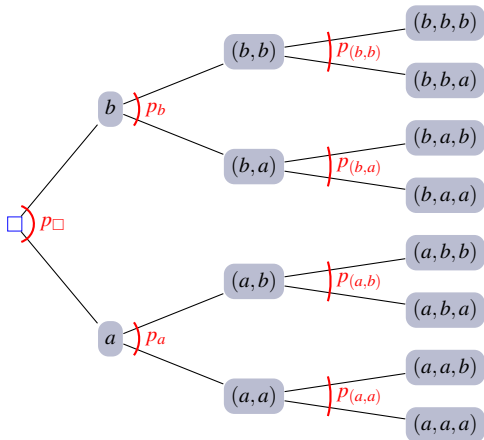
$$s = (x_1, x_2, \dots, x_n), \quad x_k \in \mathcal{X}, \quad n \geq 0.$$

The standard event tree becomes a **probability tree** by attaching to each node s a local **probability mass function** p_s on \mathcal{X} with associated **expectation operator** E_s .



Discrete-time uncertain processes

Precise probability trees



Precise probability trees

Calculating global expectations from local ones

Consider a function $g: \mathcal{X}^n \rightarrow \mathbb{R}$ of the first n variables:

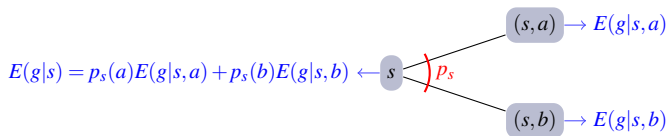
$$g = g(X_1, X_2, \dots, X_n)$$

We want to calculate its **expectation** $E(g|s)$ in $s = (x_1, \dots, x_k)$.

Theorem (Law of Iterated Expectation)

Suppose we know $E(g|s, x)$ for all $x \in \mathcal{X}$, then we can calculate $E(g|s)$ by **backwards recursion** using the local model p_s :

$$E(g|s) = \underbrace{E_s}_{\text{local}}(E(g|s, \cdot)) = \sum_{x \in \mathcal{X}} p_s(x) E(g|s, x).$$



Precise probability trees

Calculating global expectations from local ones

All expectations $E(g|x_1, \dots, x_k)$ in the tree can be calculated from the local models as follows:

- 1 start in the final cut \mathcal{X}^n and let:

$$E(g|x_1, x_2, \dots, x_n) = g(x_1, x_2, \dots, x_n);$$

- 2 do backwards recursion using the Law of Iterated Expectation:

$$E(g|x_1, \dots, x_k) = \underbrace{E_{(x_1, \dots, x_k)}}_{\text{local}}(E(g|x_1, \dots, x_k, \cdot))$$

- 3 go on until you get to the root node \square , where:

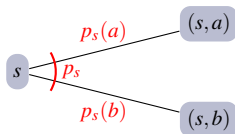
$$E(g|\square) = E(g).$$

Imprecise probability models

Sets of mass functions

Major restrictive assumption

Until now, we have assumed that we have **sufficient information** in order to specify, in each node s , a probability mass function p_s on the set \mathcal{X} of possible values for the next state.



More general uncertainty models

We consider **credal sets** as more general uncertainty models: **closed convex subsets** of $\Sigma_{\mathcal{X}}$.

Imprecise probability trees

Definition and interpretation

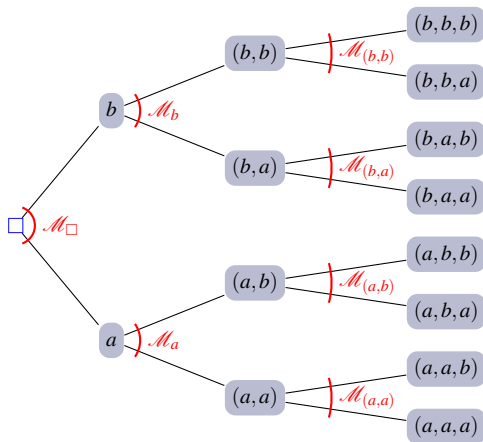
Definition

An **imprecise probability tree** is a probability tree where in each node s the local uncertainty model is an **imprecise probability model** \mathcal{M}_s , or equivalently, its associated **upper expectation** \bar{E}_s :

$$\bar{E}_s(f) = \max \{E_p(f) : p \in \mathcal{M}_s\} \text{ for all real maps } f \text{ on } \mathcal{X}.$$

Imprecise probability trees

Definition and interpretation



Imprecise probability trees

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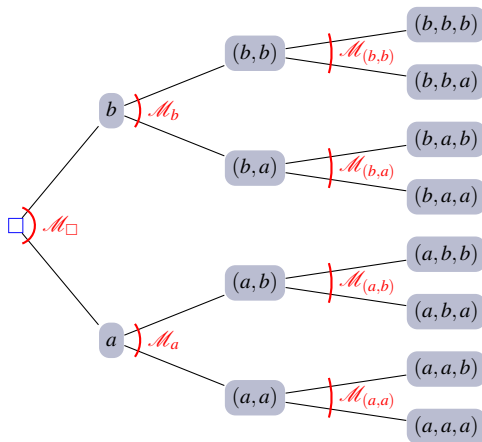
$$\bar{E}_s(f) = \max \{E_p(f) : p \in \mathcal{M}_s\} \text{ for all real maps } f \text{ on } \mathcal{X}.$$

An imprecise probability tree can be seen as an infinity of **compatible** precise probability trees: choose in each node s a probability mass function p_s from the set \mathcal{M}_s .



Imprecise probability trees

Definition and interpretation



Imprecise probability trees

Associated lower and upper expectations

For each real map $g = g(X_1, \dots, X_n)$, each node $s = (x_1, \dots, x_k)$, and each such **compatible precise probability tree**, we can calculate the expectation

$$E(g|x_1, \dots, x_k)$$

using the backwards recursion method described before.

By varying over each compatible probability tree, we get a **closed real interval**:

$$[\underline{E}(g|x_1, \dots, x_k), \bar{E}(g|x_1, \dots, x_k)]$$

We want a better, more efficient method to calculate these **lower** and **upper expectations** $\underline{E}(g|x_1, \dots, x_k)$ and $\bar{E}(g|x_1, \dots, x_k)$.

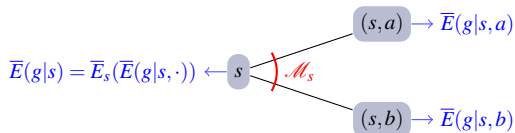
Imprecise probability trees

The Law of Iterated Expectation

Theorem (Law of Iterated Expectation)

Suppose we know $\bar{E}(g|s,x)$ for all $x \in \mathcal{X}$, then we can calculate $\bar{E}(g|s)$ by *backwards recursion* using the local model \bar{E}_s :

$$\bar{E}(g|s) = \underbrace{\bar{E}_s}_{\text{local}}(\bar{E}(g|s, \cdot)) = \max_{p_s \in \mathcal{M}_s} \sum_{x \in \mathcal{X}} p_s(x) \bar{E}(g|s, x).$$



The complexity of calculating the $\bar{E}(g|s)$, as a function of n , is therefore essentially the same as in the precise case!

Precise Markov chains

Definition

Definition

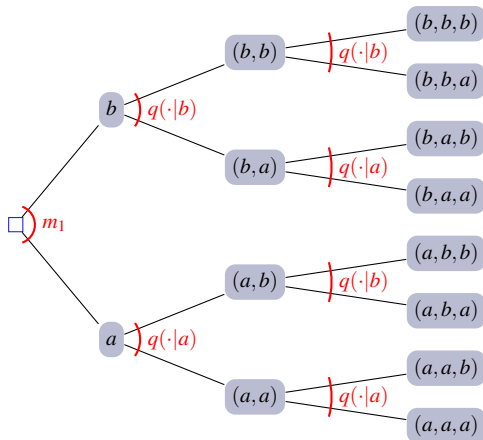
The uncertain process is a stationary **precise Markov chain** when all \mathcal{M}_s are singletons (precise), and

- 1 $\mathcal{M}_\square = \{m_1\}$,
- 2 the **Markov Condition** is satisfied:

$$\mathcal{M}_{(x_1, \dots, x_n)} = \{q(\cdot | x_n)\}.$$

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- 2 the **Markov Condition** is satisfied:

$$\mathcal{M}_{(x_1, \dots, x_n)} = \{q(\cdot | x_n)\}.$$

For each $x \in \mathcal{X}$, the transition mass function $q(\cdot | x)$ corresponds to an expectation operator:

$$E(f|x) = \sum_{z \in \mathcal{X}} q(z|x)f(z).$$

Precise Markov chains

Transition operators

Definition

Consider the linear transformation T of $\mathcal{L}(\mathcal{X})$, called **transition operator**:

$$T: \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{X}): f \mapsto Tf$$

where Tf is the real map given by, for any $x \in \mathcal{X}$:

$$Tf(x) := E(f|x) = \sum_{z \in \mathcal{X}} q(z|x)f(z)$$

T is the dual of the linear transformation with **Markov matrix** M , with elements $M_{xy} := q(y|x)$.

Precise Markov chains

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Then the **Law of Iterated Expectation** yields:

$$E_n(f) = E_1(T^{n-1}f), \text{ and dually, } m_n^T = m_1^T M^{n-1}.$$

Complexity is linear in the number of time steps n .

Imprecise Markov chains

Definition

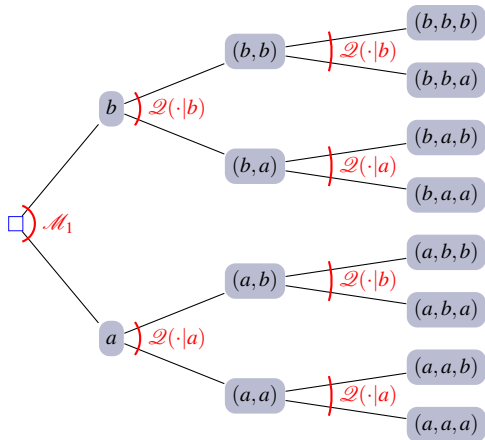
Definition

The uncertain process is a stationary imprecise Markov chain when the **Markov Condition** is satisfied:

$$\mathcal{M}_{(x_1, \dots, x_n)} = \mathcal{Q}(\cdot | x_n).$$

Imprecise Markov chains

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$$\mathcal{M}_{(x_1, \dots, x_n)} = \mathcal{Q}(\cdot | x_n).$$

An imprecise Markov chain can be seen as an infinity of probability trees.

For each $x \in \mathcal{X}$, the local transition model $\mathcal{Q}(\cdot | x)$ corresponds to **lower** and **upper expectation operators**:

$$\begin{aligned} \underline{E}(f|x) &= \min \{ E_p(f) : p \in \mathcal{Q}(\cdot|x) \} \\ \bar{E}(f|x) &= \max \{ E_p(f) : p \in \mathcal{Q}(\cdot|x) \}. \end{aligned}$$



Imprecise Markov chains

Lower and upper transition operators

Definition

Consider the **non-linear** transformations \underline{T} and \bar{T} of $\mathcal{L}(\mathcal{X})$, called **lower** and **upper transition operators**:

$$\underline{T}: \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{X}): f \mapsto \underline{T}f$$

$$\bar{T}: \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{X}): f \mapsto \bar{T}f$$

where the real maps $\underline{T}f$ and $\bar{T}f$ are given by:

$$\underline{T}f(x) := \underline{E}(f|x) = \min \{E_p(f) : p \in \mathcal{Q}(\cdot|x)\}$$

$$\bar{T}f(x) := \bar{E}(f|x) = \max \{E_p(f) : p \in \mathcal{Q}(\cdot|x)\}$$

Imprecise Markov chains

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Then the **Law of Iterated Expectation** yields:

$$\underline{E}_n(f) = \underline{E}_1(\underline{T}^{n-1}f) \text{ and } \bar{E}_n(f) = \bar{E}_1(\bar{T}^{n-1}f).$$

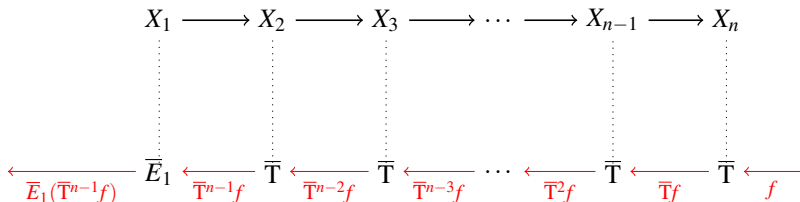
Complexity is still linear in the number of time steps n .

Imprecise Markov chains

Message passing

Important observation

The backpropagation can be seen as **message passing**.



A special credal network

under epistemic irrelevance

An imprecise Markov chain can also be depicted as follows:

$$X_1 \longrightarrow X_2 \longrightarrow X_3 \longrightarrow \dots \longrightarrow X_{n-1} \longrightarrow X_n$$

Interpretation of the graph

Conditional on X_k we have that X_1, \dots, X_{k-1} are **epistemically irrelevant** to X_{k+1}, \dots, X_n :

$$\bar{E}(f(X_{k+1}, \dots, X_n) | X_1, \dots, X_{k-1}, X_k) = \bar{E}(f(X_{k+1}, \dots, X_n) | X_k)$$

Credal networks under epistemic irrelevance

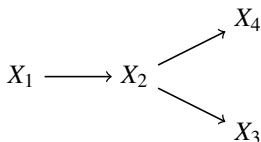
Definition

The graphical structure is interpreted as follows:

Conditional on the parents, the non-parent non-descendants of each node are **epistemically irrelevant** to it.

Credal networks under epistemic irrelevance

Example



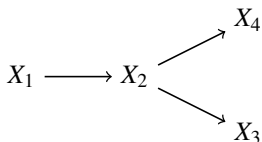
- X_1 is epistemically irrelevant to X_3 , conditional on X_2
- X_3 **need not be** epistemically irrelevant to X_1 , conditional on X_2 .

Conclusion

X_1 and X_3 **need not be** epistemically, and certainly not strongly independent, conditional on X_2 .

Credal networks under epistemic irrelevance

Example



- X_3 is epistemically irrelevant to X_4 , conditional on X_2
- X_4 is epistemically irrelevant to X_3 , conditional on X_2 .

Conclusion

X_3 and X_4 are **epistemically**, but not necessarily strongly, **independent**, conditional on X_2 .

Credal networks under epistemic irrelevance

Some separation properties



Figure: I_2 separates T from I_1 .

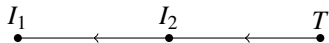


Figure: I_2 doesn't separate T from I_1 .

Conclusion

For a variable T to be separated from I_2 by a variable I_1 , **arrows should point from I_2 to T** .

Credal networks under epistemic irrelevance

As an expert system

Message passing algorithm

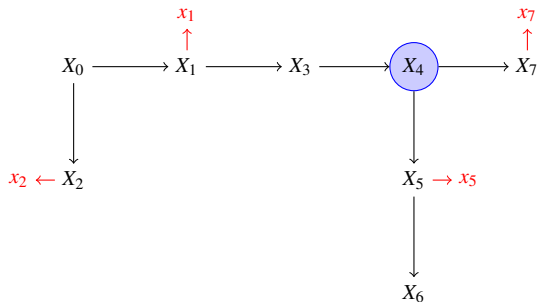
- when the credal network is a (Markov) **tree**
- treated as an expert system
- **linear complexity** in the number of nodes

Python code

- written by Filip Hermans
- testing and connection with strong independence by Alessandro Antonucci

An example

A particular Markov tree

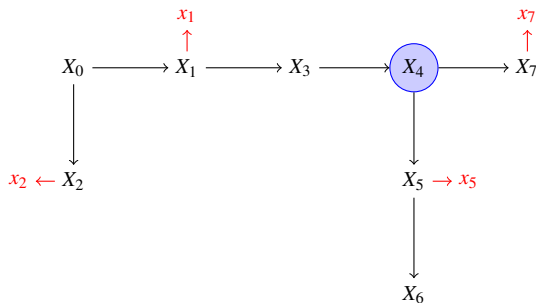


We are looking for:

$$\underline{E}(f(X_4) | x_1, x_2, x_5, x_7)$$

An example

A particular Markov tree

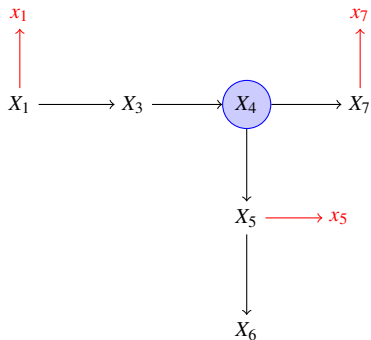


This is the unique μ such that:

$$\underline{E}([f(X_4) - \mu]I_{\{x_1\}}I_{\{x_2\}}I_{\{x_5\}}I_{\{x_7\}}) = 0$$

An example

A particular Markov tree

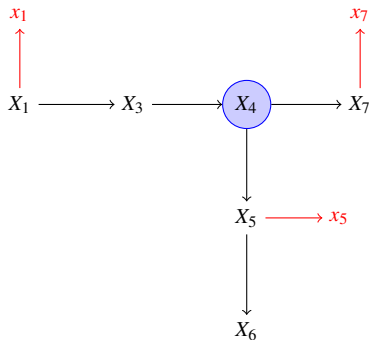


This is the unique μ such that:

$$\underline{E}([f(X_4) - \mu]I_{\{x_5\}}I_{\{x_7\}}|x_1) = 0$$

An example

A particular Markov tree

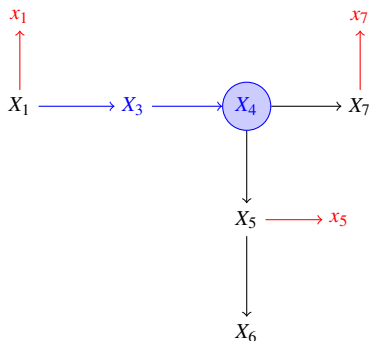


This is the unique μ such that:

$$\underline{E}_3(\underline{E}_4([f(X_4) - \mu] \underline{E}_5(\{x_5\} | X_4) \underline{E}_7(\{x_7\} | X_4) | X_3) | x_1) = 0$$

An example

A particular Markov tree



This is the unique μ such that:

$$\underline{E}_3 \left(\underbrace{\underline{E}_4 \left([f(X_4) - \mu] \overbrace{\underline{E}_5(\{x_5\} | X_4) \underline{E}_7(\{x_7\} | X_4)}^{\text{passed to the backbone}} \right) | X_3 \right) | x_1 \right) = 0$$

passed along the backbone

Literature



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