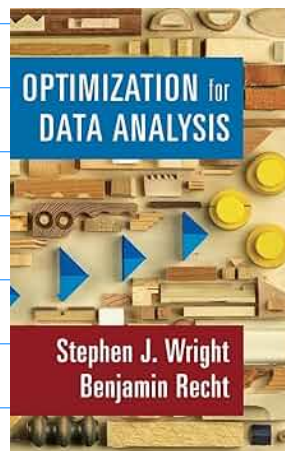
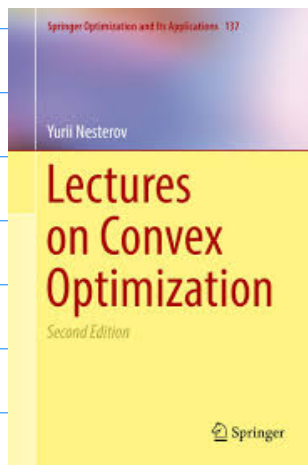


Complexity in continuous optimization (2/6)

February 6, 2024

Today: Accelerated gradient from convex to nonconvex

References:



“Convex Until Proven Guilty”: Dimension-Free Acceleration of Gradient Descent on Non-Convex Functions

ICML 2017

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Mathematical Programming (2021) 185:315–355
<https://doi.org/10.1007/s10107-019-01431-x>

FULL LENGTH PAPER

Series A

Lower bounds for finding stationary points II: first-order methods

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Where we stand:

minimize $f(x)$, $f \in C_{L}^{1,1}$ (∇f exists at every $x \in \mathbb{R}^n$
 $x \in \mathbb{R}^n$ and $\|\nabla f(x) - \nabla f(y)\| \leq L\|x-y\|$)

Goal: Bound the number of iterations/gradient evaluations/function evaluations to find a point such that $\|\nabla f(x)\| \leq \varepsilon$

One algorithm: Gradient descent

$$x_{k+1} = x_k - \alpha_k \nabla f(x_k) , \alpha_k > 0 \text{ (e.g. } \alpha_k = \frac{1}{L} \text{)}$$

Complexity: $\|\nabla f(x_k)\| \leq \varepsilon$ after at most $O(\varepsilon^{-2})$ iterations
(Valid $\forall \varepsilon > 0$, implies $\liminf_{k \rightarrow \infty} \|\nabla f(x_k)\| = 0$)
Upper bound is $O(\varepsilon^{-2})$

Lower bound matches the upper bound: There exists a function $f \in C_{L}^{1,1}$ such that GD takes exactly ε^{-2} iterations to reach a point such that $\|\nabla f(x)\| \leq \varepsilon$.

\Rightarrow Sharp analysis (Lower and upper bound match)

\Rightarrow How can we improve these results? Look at subclasses of $C_{L}^{1,1}$ functions

Today:

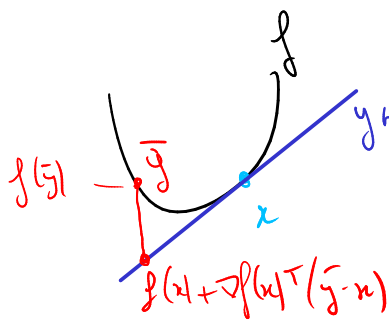
- Convex
- $C_{L}^{1,1} + C^2$ (nonconvex)

① Convex optimization

In this part, we suppose that f is $C_{L}^{1,1}$ and convex or strongly convex.

$f \in C_{L}^{1,1}$ is convex $\Leftrightarrow \forall (x, y) \in (\mathbb{R}^n)^2$,

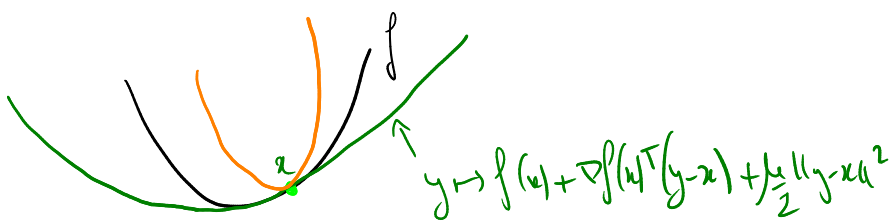
$$f(y) \geq f(x) + \underbrace{\nabla f(x)^T (y-x)}_{\text{linear function of } y}$$



Property: If f is C^1 convex, then $\nabla f(\bar{x}) = 0 \Leftrightarrow \bar{x} \in \text{argmin}_x f(x)$
 (\bar{x} global minimum of f)

$f \in C_{L}^{1,1}$ is μ -strongly convex $\Leftrightarrow \forall (x, y) \in (\mathbb{R}^n)^2$,

$$f(y) \geq f(x) + \nabla f(x)^T (y-x) + \underbrace{\frac{\mu}{2} \|y-x\|^2}_{\text{quadratic function of } y}$$



(NB: $f \in C_{L}^{1,1} \Rightarrow \forall (x, y) \in (\mathbb{R}^n)^2$, $f(y) \leq f(x) + \nabla f(x)^T (y-x) + \frac{L}{2} \|y-x\|^2$)

Property: If $f \in C_{L}^{1,1}$ and μ -strongly convex, then it has a unique minimum which is the unique solution of $\nabla f(x) = 0_{\mathbb{R}^n}$

f is μ -strongly convex $\Leftrightarrow x \mapsto \underbrace{f(x) - \frac{\mu}{2} \|x\|^2}_{\bar{f}}$ is convex

minimize $\bar{f}(y) - \bar{f}(x) - \nabla \bar{f}(x)^T (y-x)$
 x, y

$$f\left(\begin{bmatrix} x_1 \\ x_2 \end{bmatrix}\right) = x_1^2$$

\hookrightarrow Gradient descent on convex problems

$$\rightarrow x_{k+1} = x_k - \alpha_k \nabla f(x_k) = x_k - \frac{1}{L} \nabla f(x_k)$$

$\rightarrow f \in C_{L}^{1,1}$ and convex, and has at least 1 minimum

Then, GD reaches a point x_k such that

$$f(x_k) - \min_{x \in \mathbb{R}^d} f(x) \leq \varepsilon \quad \text{in at most } O(\varepsilon^{-1}) \text{ iterations}$$

$$\Leftrightarrow \|\nabla f(x_k)\| \leq \varepsilon \quad \text{in at most } O(\varepsilon^{-1}) \text{ iterations}$$

Typical criterion for complexity in the convex setting

(replaces $\|\nabla f(x)\| \leq \varepsilon$ used in the nonconvex setting)

\Rightarrow Indicates how far the current function value is from the minimum value

Proof: • Since f is $C_{L}^{1,1}$, $\forall k \in \mathbb{N}$,

$$\begin{aligned} f\left(x_k - \frac{1}{L} \nabla f(x_k)\right) &\leq f(x_k) + \nabla f(x_k)^T \left(x_k - \frac{1}{L} \nabla f(x_k) - x_k\right) + \frac{L}{2} \left\|x_k - \frac{1}{L} \nabla f(x_k) - x_k\right\|^2 \\ &\leq f(x_k) - \frac{1}{2L} \|\nabla f(x_k)\|^2 \end{aligned}$$

x_{k+1}

$$f(x_{k+1}) \leq f(x_k) - \frac{1}{2L} \|\nabla f(x_k)\|^2$$

Let $x^* \in \operatorname{argmin}_x f(x)$ and $f^* = f(x^*) = \min_x f(x)$

By convexity, $f(x^*) \geq f(x_k) + \nabla f(x_k)^T (x^* - x_k)$

$$\Leftrightarrow f(x_k) \leq f(x^*) + \nabla f(x_k)^T (x_k - x^*)$$

Hence, $f(x_{k+1}) \leq f(x^*) + \nabla f(x_k)^T (x_k - x^*) - \frac{1}{2L} \|\nabla f(x_k)\|^2$

Then $f(x_{k+1}) - f(x^*) \leq \nabla f(x_k)^T (x_k - x^*) - \frac{1}{2L} \|\nabla f(x_k)\|^2$

$$= \frac{L}{2} \left(\|x_k - x^*\|^2 - \|x_{k+1} - x^*\|^2 \right)$$

$$\uparrow$$

$$x_k - \frac{1}{L} \nabla f(x_k)$$

Suppose that $f(x_k) - f(x^*) > \varepsilon \quad \forall k = 0, \dots, K$

$$\begin{aligned} \text{Then} \quad \sum_{k=0}^{K-1} f(x_{k+1}) - f(x^*) &\leq \frac{L}{2} \sum_{k=0}^{K-1} (\|x_k - x^*\|^2 - \|x_{k+1} - x^*\|^2) \\ &= \frac{L}{2} (\|x_0 - x^*\|^2 - \|x_K - x^*\|^2) \end{aligned}$$

$$K \varepsilon < \underbrace{\sum_{k=0}^{K-1} (f(x_{k+1}) - f(x^*))}_{> \varepsilon} \leq \frac{L}{2} \|x_0 - x^*\|^2$$

$$\Rightarrow K < \frac{L}{2} \|x_0 - x^*\|^2 \varepsilon^{-1} = O(\varepsilon^{-1})$$

For μ -strongly convex case, can show

$$f(x_k) - f(x^*) \leq \left(1 - \frac{\mu}{L}\right)^k (f(x_0) - f(x^*))$$

(Behind this: $\|\nabla f(x)\|^2 \geq 2\mu (f(x) - f(x^*))$, natural consequence of strong convexity inequality)

$$f \stackrel{C^{1,1}}{L} \Rightarrow f(x_{k+1}) \leq f(x_k) - \frac{1}{2L} \|\nabla f(x_k)\|^2$$

$$\leq f(x_k) - \frac{\mu}{L} (f(x_k) - f(x^*))$$

$$f(x_{k+1}) - f(x^*) \leq f(x_k) - f(x^*) - \frac{\mu}{L} (f(x_k) - f(x^*)) = \left(1 - \frac{\mu}{L}\right) (f(x_k) - f(x^*))$$

→ The proof in the strongly convex setting is (in some way) easier than the proof in the convex setting

Complexity of GD : $O\left(\frac{L}{\mu} \ln(\varepsilon^{-1})\right)$

Same GD algorithm	$f \stackrel{C^{1,1}}{L}$
f non-convex	$\ \nabla f(x)\ \leq \varepsilon \quad O(\varepsilon^{-2})$
f convex	$f(x) - \min_x f(x) \leq \varepsilon \quad O(\varepsilon^{-1})$
f μ -strongly convex	$\quad \quad \quad \quad \quad \quad \quad \quad O\left(\frac{L}{\mu} \ln(\varepsilon^{-1})\right)$

↑ constant in term of ε but important for the complexity
 L/μ : "condition number"

Pb: In the convex/strongly convex settings, the upper bounds for GD do not match the lower bounds:

Abramovskii & Yudin 1983

Existence result:
 there exists an algorithm that uses only one gradient per iteration

→ If $f \in C_{L, \mu}^{1,1}$ convex, there exists an algorithm with complexity $O(\varepsilon^{-1/2})$

→ If $f \in C_{L, \mu}^{1,1}$ μ -strongly convex, \exists algorithm with complexity $O\left(\sqrt{\frac{L}{\mu}} \ln(\varepsilon^{-1})\right)$

$\mu \leq L \implies \sqrt{\frac{L}{\mu}} \leq \frac{L}{\mu}$

→ Algorithm was unknown until Yurii Nesterov discovered it in 1983

Accelerated gradient / Nesterov's method ($x_0 \in \mathbb{R}^n$)

$$x_{k+1} = x_k - \alpha_k \nabla f\left(x_k + \underbrace{\beta_k (x_k - x_{k-1})}_{\text{momentum term}}\right) + \beta_k (x_k - x_{k-1})$$

$$\alpha_k > 0, \beta_k > 0 \quad (\beta_0 = 0, x_{-1} = x_0)$$

Equivalent formulation:

$$\begin{cases} y_k = x_k + \beta_k (x_k - x_{k-1}) \\ x_{k+1} = y_k - \nabla f(y_k) \end{cases}$$

$$\begin{aligned} \beta_0 &= 0 \\ x_{-1} &= x_0 \\ y_0 &= x_0 \end{aligned}$$

Note: Momentum methods in RL (SGD with momentum, Adam)

$$x_{k+1} = x_k - \alpha_k \nabla f_{ii}(x_k) + \beta_k (x_k - x_{k-1})$$

$\nabla f(x_k)$: stochastic gradient \approx sample

Analyzing Nesterov's method

$$\begin{cases} y_k = x_k + \beta_k (x_k - x_{k-1}) \\ x_{k+1} = y_k - \alpha_k \nabla f(y_k) \end{cases} \quad y_0 = x_0, \beta_0 = 0, x_{-1} = x_0$$

a) f is μ -strongly convex

$$\rightarrow \alpha_k = \frac{1}{L}, \quad \beta_k = \frac{\sqrt{k} - 1}{\sqrt{k} + 1} \quad \text{where } \kappa = \frac{L}{\mu}$$

Key difference with GD analysis

$$f^* = \min f(x)$$

• In GD, we look at $f(x_k) - f^*$ and we show

$$f(x_k) - f^* \leq \left(1 - \frac{\mu}{L}\right)^k (f(x_0) - f^*) \\ = \left(1 - \frac{1}{\kappa}\right)^k (f(x_0) - f^*)$$

"Lyapunov function for GD"

• In AG (Accelerated Gradient)'s analysis, we use

$$V_k := f(x_k) - f^* + \frac{L}{2} \|x_k - x^* - \rho^2 (x_{k-1} - x^*)\|^2$$

with $\rho^2 = \left(1 - \frac{1}{\sqrt{k}}\right) = 1 - \sqrt{\frac{\mu}{L}}$

\Rightarrow the analysis shows $V_k \leq \rho^k V_0$

$$V_0 = f(x_0) - f^* + \frac{L}{2} \|(1-\rho^2)(x_0 - x^*)\|^2$$

$$= f(x_0) - f^* + \frac{\mu}{2} \|x_0 - x^*\|^2$$

\Rightarrow Gives the $O\left(\sqrt{\frac{L}{\mu}} \ln(\varepsilon^{-1})\right)$ bound

vs $O\left(\frac{L}{\mu} \ln(\varepsilon^{-1})\right)$ for gradient descent

b) Convex case (Not strongly convex)

Nesterov's implementation of accelerated gradient

$$\begin{cases} y_k = x_k + \beta_k (x_k - x_{k-1}) \\ x_{k+1} = y_k - \frac{1}{L} \nabla f(y_k) \end{cases}$$

$$\begin{aligned} \beta_0 &= 0 \\ \beta_k &= \rho_k \rho_{k-1}^2 \quad k \geq 1 \\ \rho_0 &= \rho_{-1} = 0 \end{aligned}$$

Not most intuitive part of the algorithm: $\{\rho_k\}$ and $\{\beta_k\}$ is defined independently of the problem (f) and x_0

ρ_{k+1} is the positive root of

$$1 + \rho_{k+1} (\rho_k^2 - 1) - \rho_{k+1}^2 = 0$$

$$\rho_{k+1} = \frac{1 + \sqrt{1 + 4\rho_k^2}}{2}$$

\rightarrow Analysis relies on the Lyapunov function

$$W_k = f(x_k) - f^* + \frac{L}{2} \|(x_k - x^*) - \rho_{k-1}^2 (x_{k-1} - x^*)\|^2$$

\rightarrow show $W_k \leq \rho_{k-1}^2 \dots \rho_1^2 W_1 = (1 - \rho_{k-1}^2)^2 W_1$

and $W_1 \leq \frac{L}{2} \|x_0 - x^*\|^2$

$$\rightarrow 1 - \rho_k^2 \leq \frac{2}{k+2} \quad \forall k \text{ (by definition of } \rho_k)$$

Overall, prove $f(x_k) - f^* \leq V_k \leq \frac{2L}{(k+1)^2} \|x_0 - x^*\|^2$

$\Rightarrow f(x_k) - f^* \leq \epsilon$ after at most $O(\epsilon^{-1/2})$ iterations.

\rightarrow Nesterov's method attains the upper bound $O(\epsilon^{-1/2})$

\rightarrow Nesterov also showed a lower bound for the convex setting

$$f(x) = \frac{1}{2} x^T A x - e_1^T x \quad \text{convex quadratic}$$

$$= \frac{1}{2} (x_1 - 1)^2 + \frac{1}{2} \sum_{i=2}^{m-1} (x_i - x_{i+1})^2$$

$$A = \begin{bmatrix} 2 & -1 & & & 0 \\ -1 & 2 & & & \\ & & \ddots & \ddots & \\ 0 & & & -1 & 2 \\ & & & & & & -1 & 2 \end{bmatrix}, e_1 = \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

unique solution

$$x^* = \begin{bmatrix} 1 - \frac{1}{m+1} \\ \vdots \\ 1 - \frac{m}{m+1} \end{bmatrix}$$

Any algorithm starting from $x_0 = 0$ and using only 1 gradient per iteration will satisfy

$$f(x_k) - f^* \geq \frac{3}{8(k+1)^2} \|x_0 - x^*\|^2 \quad \forall k \leq \frac{m}{2} - 1$$

Nesterov's example: dimension " $\frac{1}{\epsilon}^{-1/2}$ " but quadratic convex

Carr, Gold, Jost examples: dimension 1 or 2, but very nonlinear nonquadratic

② Acceleration in the nonconvex setting

→ If f is just $C_{L}^{1,1}$ and nonconvex, cannot do better than GD!

→ If f is $C_{L}^{1,1}$, nonconvex and $C^{2,2}$, then you can do better than GD, and you can do that with a variant of Nesterov's method (Cannons et al '17)

Idea: "Convex until proven guilty"

- Run AG as if the function were strongly convex
- Two cases:
 - Either the method works as in the strongly convex setting
 - Or it doesn't, and take a negative curvature step instead of doing 1 iteration of AG

Negative curvature

f nonconvex C^2 :

$$\bar{x} \text{ (argmin)}_x f(x) \Rightarrow \begin{cases} \nabla f(\bar{x}) = 0 \\ \nabla^2 f(\bar{x}) \succeq 0 \end{cases}$$

$$\forall v \in \mathbb{R}^n, \quad \uparrow \quad v^\top \nabla^2 f(\bar{x}) v \geq 0$$

When f is convex C^2 , $\nabla^2 f(x) \succeq 0$ true $\forall x$!

For nonconvex f , if $\nabla^2 f(x) \succeq 0$ does not hold,
the function decreases from x along any

direction v such that $v^T \nabla^2 f(x) v < 0$

\uparrow
Negative curvature direction (and $v^T \nabla f(x) \leq 0$)

Negative curvature step: $x + \alpha v$
 $\alpha > 0$

Algorithm $(x_0, L, \epsilon, \epsilon_H)$

L Lipschitz constant for the gradient

$\epsilon > 0$ for complexity: used in the algorithm

For $k=0, 1, \dots$

\rightarrow Define $\hat{f}_k(z) = f(z) + \epsilon_H \|z - x_k\|^2 \approx$ Lyapunov function

\rightarrow Run AG on \hat{f}_k to find a point z_k such that $\|\nabla \hat{f}_k(z_k)\| \leq \epsilon/10$

assuming that the function \hat{f}_k is ϵ_H -strongly convex

\downarrow
true if f is convex, false otherwise

At every iteration of AG, check whether the strong convexity inequality holds between every pair of points that it computed, and stop if the inequality is violated.

$$f(u) \geq \hat{f}_k(u) + \frac{\epsilon_H}{2} \|u - v\|^2$$

• If AG stops with $\|\nabla \hat{f}_k(x_k)\| \leq \frac{\epsilon}{L}$,
 define $x_{k+1} = x_k$.

The function
 is μ -strongly
 convex \rightarrow
 "enough" around
 x_k for AG
 to fail

• Otherwise, get a pair (u_k, v_k) that
 violates the strong convexity inequality.
 and use $u_k - v_k$ as a negative
 curvature direction.

$$x_{k+1} = x_k + \eta (u_k - v_k)$$

Analysis: At every iteration k ,

$$f(x_k) - f(x_{k+1}) \geq O\left(\frac{\|\nabla f(x_k)\|^2}{\epsilon_k}\right) \geq O\left(\frac{\epsilon^2}{\epsilon_k}\right)$$

if $\|\nabla f(x_k)\| \geq \epsilon$

• $f(x_k) - f(x_{k+1}) \geq O(\epsilon_k^3)$
 if AG fails

\Rightarrow The method computes x_k s.t. that $\|\nabla f(x_k)\| \leq \epsilon$
 in at most $O(\max(\epsilon^{-2} \epsilon_k, \epsilon_k^{-3}))$ iterations

\Rightarrow With $\epsilon_k = \epsilon^{1/2}$, get $O(\epsilon^{-3/2})$ iterations

Better than $O(\epsilon^{-2})$ for GD!

→ Every call to AG terminates (with success/failure) after at most $O\left(\frac{\sqrt{L+2\varepsilon_H}}{\varepsilon_H} \ln\left(\frac{1}{\varepsilon}\right)\right) = O\left(\sqrt{\frac{1}{\varepsilon_H}} \ln\left(\frac{1}{\varepsilon}\right)\right)$

⇒ Total number of iterations:

$$O\left(\frac{1}{\sqrt{\varepsilon_H}} \ln\left(\frac{1}{\varepsilon}\right)\right) \times O\left(\max(\varepsilon^{-2} \varepsilon_H, \varepsilon_H^{-3})\right)$$

$$\varepsilon_H = \varepsilon^{1/2}$$

$$\hookrightarrow O\left(\varepsilon^{-7/4} \ln(1/\varepsilon)\right) = \tilde{O}\left(\varepsilon^{-7/4}\right)$$

⇒ still better (for small ε) than $O(\varepsilon^{-2})$

$\varepsilon^{-7/4}$: Best known upper bound for algorithms that use only gradient information to optimize a $C^{1,1} \cap C^{2,2}$ function

Lower bound (for this class of methods/functions): $O\left(\varepsilon^{-12/7}\right)$

⇒ Achieved for an example with dimension $O(\varepsilon^{-12/7})$

Takeaways

- AG better than GD for convex (strongly convex) versions
- Analyser: Lyapunov functions + convexity inequalities
- Extension to nonconvex (+negative curvature)