

Reducing graph coloring to stable set without symmetry

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Coloring and perfect 0-1 matrices

- $(P) = \{Ax \leq b\}$ is *TDI*:

$$\begin{aligned} & \min\{y^\top b : y^\top A \geq w^\top\} \\ = & \min\{y^\top b : y^\top A \geq w^\top, y \text{ integer}\} \quad (\forall w \text{ integer}) \end{aligned}$$

- $P = \{x : Ax \leq b\}$ is *integer*:

$$\begin{aligned} & \max\{w^\top x : x \in P\} \\ = & \max\{w^\top x : x \in P, x \text{ integer}\} \quad (\forall w) \end{aligned}$$

- A is 0-1: the clique-matrix $\{\text{maximal cliques}\} \times \{\text{vertices}\}$

$$\begin{aligned} \alpha(G) := & \max\{\mathbf{1}^\top x : Ax \leq \mathbf{1}, x \geq \mathbf{0}, x \text{ integer}\} =: \omega(\overline{G}) \\ \overline{\chi}(G) := & \min\{y^\top \mathbf{1} : y^\top A \geq \mathbf{1}^\top, y \geq \mathbf{0}, y \text{ integer}\} =: \chi(\overline{G}) \end{aligned}$$

Compact LP formulation

$$\begin{aligned}\chi(G) = \min \quad & \sum_i x_i \\ \text{s.t.} \quad & x_i \geq x_i^v \quad (\forall i, \forall v) \\ & \sum_i x_i^v = 1 \quad (\forall v) \\ & x_i^u + x_i^v \leq 1 \quad (\forall i, \forall uv) \\ & x \geq \mathbf{0} \\ & x \text{ integer}\end{aligned}$$

where $G = (V, E)$ is a graph, $i \in \{1, \dots, |V|\}$, $v \in V$, $uv \in E$.

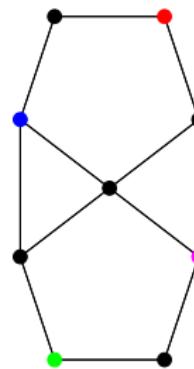
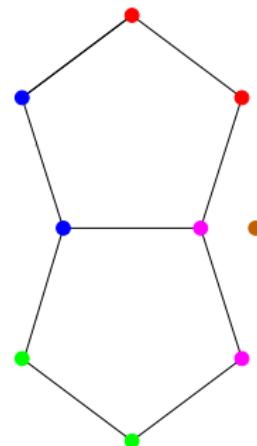
$\hat{G} = G_1 \square \dots \square G_{|V(G)|}$ add $u_1, \dots, u_{|V(G)|}$ universal

$$\chi(G) + \alpha(\hat{G}) = 2|V(G)|$$

Gallai identities

For all G (without isolated vertex)

$$\rho(G) + \nu(G) = |V(G)| = \alpha(G) + \tau(G)$$



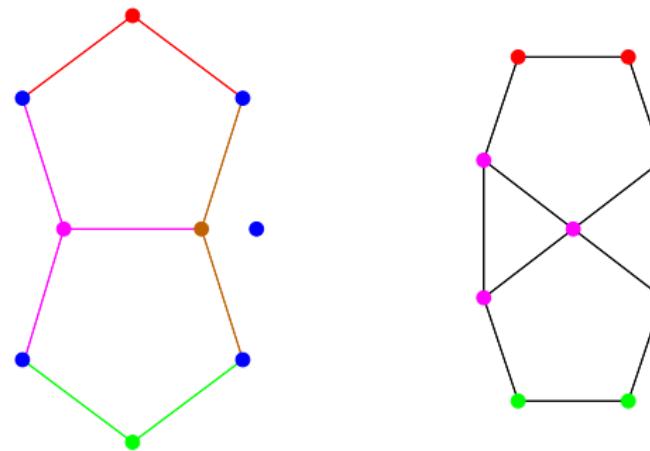
If G has no 3-cycle

$$\bar{\chi}(G) + \alpha(L(G)) = |V(G)|$$

Gallai identities

For all G (without isolated vertex)

$$\rho(G) + \nu(G) = |V(G)| = \alpha(G) + \tau(G)$$



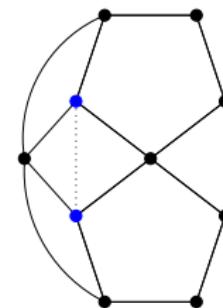
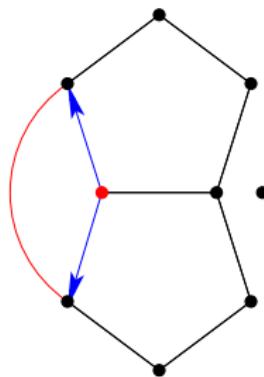
If G has no 3-cycle

$$\chi(G) + \alpha(L(G)) = |V(G)| = \alpha(G) + \overline{\chi}(L(G))$$

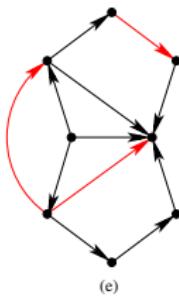
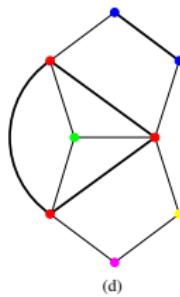
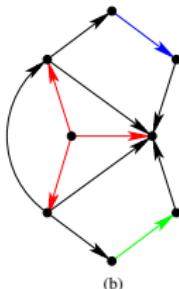
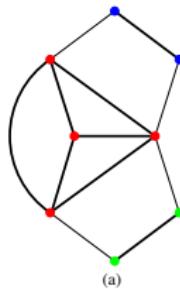
The “sandwich” line-graphs $S(G)$ of G

Choose a orientation \vec{G} of G without 3-dicycle (cliques are acyclic)

$S(\vec{G})$ is the line-graph $L(G)$ of G minus the edges corresponding to simplicial pairs of arcs



Chromatic Gallai identities



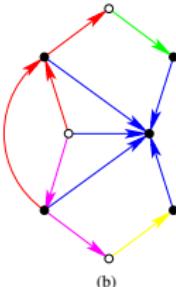
\vec{G} is an orientation of G , without 3-dicycle

$$\chi(G) + \alpha(S(\vec{G})) = |V(G)|$$

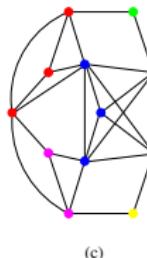
Chromatic Gallai identities



(a)



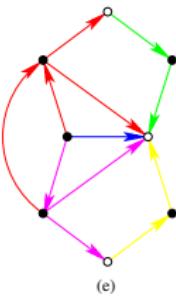
(b)



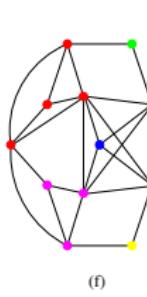
(c)



(d)



(e)



(f)

\vec{G} is an orientation of G , without 3-dicycle

$$\overline{\chi}(G) + \alpha(S(\vec{G})) = |V(G)| = \alpha(G) + \overline{\chi}(S(\vec{G}))$$

The operator Φ

Let $\beta(\cdot)$ be a graph parameter such that $\alpha(G) \leq \beta(G) \leq \bar{\chi}(G)$ ($\forall G$)

Choose an orientation \vec{G} of G without 3-dicycle and define

$$\Phi_\beta(G) := |V(G)| - \beta(S(\vec{G}))$$

Corollary. $\alpha(G) \leq \Phi_\beta(G) \leq \bar{\chi}(G)$

Proof.

$$|V(G)| - \bar{\chi}(S(\vec{G})) \leq |V(G)| - \beta(S(\vec{G})) \leq |V(G)| - \alpha(S(\vec{G}))$$

Fractional clique-cover (co-chromatic) number

$$\begin{aligned}\overline{\chi}_f(G) &= \left\{ \begin{array}{ll} \max & \mathbf{1}^\top x \\ \text{s.t.} & x(K) \leq 1 \quad (\forall \text{clique } K) \\ & x_v \geq 0 \quad (\forall v \in V) \end{array} \right. \\ &= \left\{ \begin{array}{ll} \min & \mathbf{1}^\top y \\ \text{s.t.} & \sum_{K \ni v} y_K = 1 \quad (\forall v) \\ & y_K \geq 0 \quad (\forall \text{clique } K) \end{array} \right.\end{aligned}$$

Lemma. y optimal for $G \Rightarrow x$ feasible for $S(\vec{G})$

$$x_{uv} := \sum_{\text{simplicial stars } S_K \ni uv} y_K$$

Lovász Theta function

$$\vartheta(G) = \left\{ \begin{array}{ll} \max & \sum_v \|x_v\|^2 \\ \text{s.t.} & \|x_o\|^2 = 1 \\ & x_o^\top x_v = \|x_v\|^2 \quad (\forall v) \\ & x_u^\top x_v = 0 \quad (\forall uv \in E) \end{array} \right.$$

$$= \left\{ \begin{array}{ll} \min & \|y_o\|^2 \\ \text{s.t.} & \|y_v\|^2 = 1 \quad (\forall v) \\ & y_o^\top y_v = 1 \quad (\forall v) \\ & y_u^\top y_v = 0 \quad (\forall uv \notin E) \end{array} \right.$$

Lemma. $x \in \mathbb{R}^{d \times m+1}$ optimal for $S(\vec{G}) \Rightarrow y \in \mathbb{R}^{nd \times n+1}$ feasible for G

$$y_{ov} := x_o - \sum_{u:uv \in E(\vec{G})} x_{uv} \text{ and } y_{vu} := \begin{cases} y_{ov} & \text{if } u = v \\ x_{uv} & \text{if } uv \in E(\vec{G}) \\ 0 & \text{otherwise} \end{cases}$$

Improving Lovász's ϑ bound for coloring

The sandwich theorem (Lovász 1979)

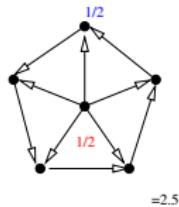
$$\alpha(G) \leq \vartheta(G) \leq \overline{\chi}_f(G) \leq \overline{\chi}(G) \quad (\forall G)$$

Theorem (C. and Meurdesoif 2014)

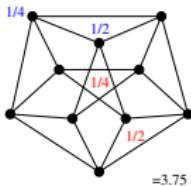
$$\alpha(G) \leq \Phi_{\overline{\chi}_f}(G) \leq \overline{\chi}_f(G) \quad \text{and} \quad \vartheta(G) \leq \Phi_{\vartheta}(G) \leq \overline{\chi}(G) \quad (\forall G)$$

	$ V $	$ E $	$\min \rho := \frac{\Phi_{\vartheta} - \vartheta}{\vartheta}$	mean ρ	max ρ
M_3	5	5	—	23.6%	—
M_4	11	35	26.8%	27.3%	28.7%
M_5	23	182	26.5%	27.6%	29.5%
M_6	47	845	26.0%	27.8%	29.5%

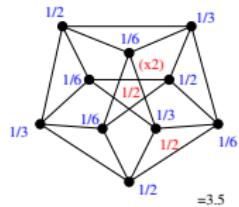
Impact of orientation



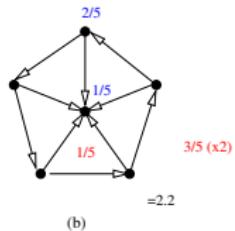
(max)
(min)



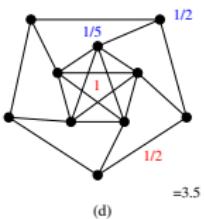
=3.75



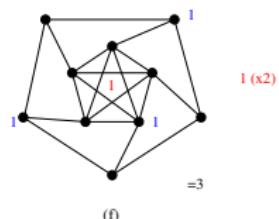
=3.5



=2.2



=3.5



1 (x2)

	2	2.2	2.25	2.5	3
6-wheel orientation (a)	α	$\bar{\chi}'_f$		$\bar{\chi}_f$	$\bar{\chi}$
orientation (b)	$\Phi_{\bar{\chi}}$		$\Phi_{\bar{\chi}'_f}$	$\Phi_{\bar{\chi}'_f}$	Φ_{α}
				$\Phi_{\bar{\chi}'_f} = \Phi_{\alpha}$	

Complexity

Theorem. (Grötschel, Lovász and Schrijver 1981)

$\vartheta(G)$ can be approximated in polynomial time (for any $\varepsilon > 0$)

Theorem. (Gvozdenović and Laurent 2008)

Computing $\overline{\chi}_f(G)$ is NP-hard

In fact, [GL08] proved that,
given any parameter $\beta(G)$ such that $\beta \in [\overline{\chi}_f, \overline{\chi}]$,
then computing $\beta(G)$ is NP-hard.

Finding α with $\Phi_{\bar{\chi}_f}$ (another proof for GL08)

$\vec{G}_k = G$ orient and add s_1, \dots, s_k universal sources

Feasible solution x for $\bar{\chi}_f(S(\vec{G}_k))$ (max)

$$x_{uv} := \begin{cases} \frac{1}{\alpha(G)} & \text{if } u = s_i \\ 0 & \text{otherwise} \end{cases}$$

Let $\beta \in [\bar{\chi}_f, \bar{\chi}]$. Start with $k := 1$ and grow it until $\Phi_\beta(\vec{G}_k) = k$.

$$\Phi_\beta(\vec{G}_k) = k \iff \alpha(G) = k$$

Proof.

$$k \leq \alpha(G) = \alpha(G_k) \leq \Phi_\beta(\vec{G}_k) \leq \Phi_{\bar{\chi}_f}(\vec{G}_k) \leq k + |V(G)| - k \frac{|V(G)|}{\alpha(G)}$$